

$(0, 1, \dots, q-1) \in GF(q)$ $\left\{ \begin{array}{l} * \\ + \end{array} \right\}$ mod q (q prime)

D): $\forall \alpha \in GF(q) \setminus \{0\} \quad \alpha^{q^1} = 1$

D): min $m: \alpha^m = 1 \rightarrow \text{ord } (\alpha) = m$

D): $\text{ord } (\alpha) = q-1 \rightarrow \alpha$ primitive elme

q -ary Hamming code $t=1$ juntatsoara

$\underline{c}^{(i)} = (0 \dots 0 \alpha 0 \dots 0) \quad \alpha \in GF(q) \setminus \{\pm 1\} \quad c_i \& \text{identifizierbar}$

$$\boxed{H \underline{c}^{\omega^T} = \underline{s}^T} \quad \leftarrow H \underline{x}^T = \underline{s}^T$$

$$\left(\begin{array}{ccc|c} : & \alpha^{(1)^T} & \dots & \alpha^{(n)^T} \\ \vdots & \vdots & \dots & \vdots \\ \alpha^{(1)^T} & \dots & \alpha^{(i)^T} & \dots & \alpha^{(n)^T} \\ \vdots & \vdots & \vdots & \vdots & \vdots \end{array} \right) \left(\begin{array}{c} 0 \\ \vdots \\ 0 \\ \alpha \\ 0 \\ \vdots \\ 0 \end{array} \right) \leftarrow \underline{s}^T = \alpha \underline{x}^T = \underline{s}^T$$

Fehlerdetektion:

- 1.) $\underline{a}^{(i)} = \underline{a}^{(j)}$ $\forall i, j = 1 \dots n \quad i \neq j$
- 2.) $\underline{a}^{(i)} \neq \underline{0}$
- 3.) $\underline{a}^{(i)}$ else neue zelle clear 1

zB: $\begin{pmatrix} 0 \\ 1 \\ 2 \\ 3 \end{pmatrix} \rightarrow \begin{pmatrix} 0 \\ \alpha \\ \vdots \\ 0 \end{pmatrix}$

$$\beta \in GF(q) \rightarrow \boxed{\underline{B} \cdot \underline{s}} \rightarrow \underline{x} = \beta^{-1} \underline{s}$$

Hamming code ist "perfekt"

$$\frac{q^{n-k}-1}{q-1} = n \quad q^{n-k} = n(q-1) + 1 = \sum_{i=0}^k \binom{n}{i} (q-1)^i$$

PERFECT KOD

(zB:) $t=1$ juntatsoara $C_4(8, 6)$ GF(2) fehler*

$$C_4(8, 6) \quad H_{2 \times 8} = \begin{pmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 0 \\ 1 & 2 & 3 & 4 & 5 & 6 & 7 & 1 \end{pmatrix}$$

$$H \cdot G^T = 0 \quad \downarrow \quad \underline{x} = -\underline{B}^T$$

G

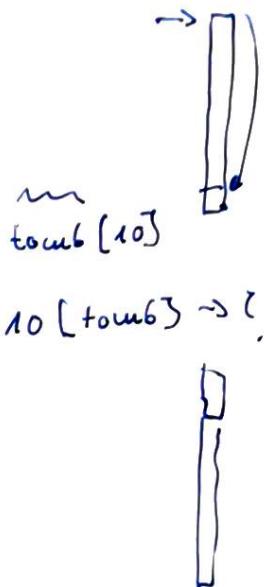
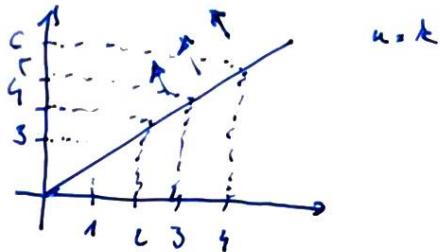
Mögigkeiten:

$$C_H(n, n-k) \quad t=1 = \left\lceil \frac{d_{min} - 1}{2} \right\rceil$$

$$3 \leq d_{min} \leq n-k+1 = 3$$

$$\downarrow \\ d_{min} = n-k+1 \text{ MDS } \heartsuit$$

$$\begin{array}{l} t=1 \\ t \geq 2 \text{ MDS} \\ \text{heute?} \end{array}$$



Polynom $GF(q)$ - Basis

$$(a_0, a_1, \dots, a_n) \in GF(q) \quad a(x) = a_0 + a_1 x + a_2 x^2 + \dots + a_n x^n$$

$$(b_0, b_1, \dots, b_n)$$

$$a_0, a_1, \dots, a_n, x \in GF(q)$$

$$\deg(a(x)) = n$$

$$b(x) = b_0 + b_1 x + b_2 x^2 + \dots + b_n x^n$$

$$c(x) = a(x) + b(x) = (a_0 + b_0) + (a_1 + b_1)x + (a_2 + b_2)x^2 + \dots + (a_n + b_n)x^n$$

$$c(x) = a(x) \cdot b(x) = a_0 b_0 + (a_1 b_0 + a_0 b_1)x + (a_2 b_0 + a_1 b_1 + a_0 b_2)x^2 + \dots + \sum_{i=0}^{\min(\deg(a(x)), \deg(b(x)))} \left(\sum_{j=0}^i a_j b_{i-j} \right) x^i$$

Polynom r.

$$c(x) = a(x) + b(x)$$

$$-c(x) = a(x) b(x)$$

Vektor r :

$$c_i := a_i + b_i \rightarrow c = a + b$$

mit $\deg(a(x))$

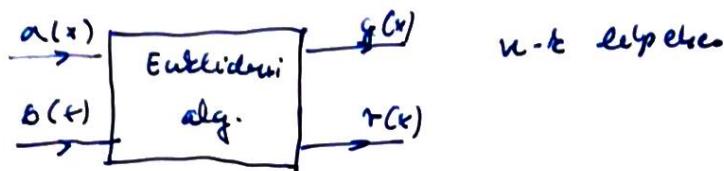
$$c_i = \sum_{j=0}^i a_j b_{i-j} \rightarrow c = a \cdot b$$

↑ konvolutiv

Polynomdivision

$$\text{Addit. } a(x), b(x) \quad \deg(a(x)) = n > \deg(b(x)) = k$$

$$a(x) = q(x)b(x) + r(x) \quad \deg(r(x)) < k$$



$n-k$ Nullstellen

(T): $a(x) \Big|_{x=x_1} = 0 \rightarrow a(x) = g(x)(x-x_1)$
 $\deg(a(x)) > \deg(g(x))$

(S): $a(x) \Big|_{x=x_1} \quad g(x)(x-x_1) \Big|_{x=x_1} + r = 0$

Köv: $a(x) = \prod_{i=1}^{\deg(a(x))} (x-x_i)$ # ggüdeste $\leq \deg(a(x))$
algebraic algorithm

Reed-Solomon Kodierung

RS Label

$$\underline{u} = (u_0, u_1, \dots, u_{n-1}) \in GF(q) \quad x_0, x_1, \dots, x_{n-1} \in GF(q) \quad x_i \neq x_j$$



$$u(x) = u_0 + u_1 x + u_2 x^2 + \dots + u_{k-1} x^{k-1}$$

q prime
 $\boxed{n = q - 1}$

$$\underline{c} = (c_0, c_1, \dots, c_{n-1})$$

$$c_0 = u(x) \Big|_{x=x_0} = u_0 + u_1 x_0 + u_2 x_0^2 + \dots + u_{k-1} x_0^{k-1}$$

$$c_1 = u(x) \Big|_{x=x_1} = u_0 + u_1 x_1 + u_2 x_1^2 + \dots + u_{k-1} x_1^{k-1}$$

$$\vdots$$

$$\underline{c}^T = \underline{u}^T \cdot G_{k \times n}$$

$$(c_0, c_1, \dots, c_{n-1})^T = (u_0, u_1, \dots, u_{k-1})^T \cdot$$

$$\left(\begin{array}{cccc} 1 & 1 & 1 & \cdots & 1 \\ x_0 & x_1 & x_2 & \cdots & x_{n-1} \\ x_0^2 & x_1^2 & x_2^2 & \cdots & x_{n-1}^2 \\ \vdots & \vdots & \ddots & & \vdots \\ x_0^{k-1} & x_1^{k-1} & x_2^{k-1} & \cdots & x_{n-1}^{k-1} \end{array} \right)$$

(1.) $d_{\min} = w_{\min} = n - \# \text{ Fehler } \leq k = w(k-1) = n-k+1$

(2.) $d_{\min} \leq n-k+1$

MOS Label

RS Ladek Koeffizienten

& primitive elem GF(q)

$$\alpha_0 = \alpha^0 = 1 \quad \alpha_1 = \alpha^1 = \alpha \quad \alpha_2 = \alpha^2, \dots, \alpha_{n-1} = \alpha^{n-1}$$

$$G_{RS} = \begin{pmatrix} 1 & 1 & 1 & \dots & 1 \\ 1 & \alpha & \alpha^2 & \dots & \alpha^{n-1} \\ 1 & \alpha^2 & \alpha^4 & \dots & \alpha^{2(n-1)} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & \alpha^{k-1} & \alpha^{2k-2} & \dots & \alpha^{(n+k)(n-1)} \end{pmatrix}$$

$$H \cdot G = 0$$

$$H \cdot G^\top = 0$$

$$H = \begin{pmatrix} 1 & 1 & 1 & \dots & \alpha^{n-1} \\ 1 & \alpha & \alpha^4 & \dots & \alpha^{2(n-1)} \\ 1 & \alpha^2 & \alpha^6 & \dots & \alpha^{3(n-1)} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & \alpha^{n-k} & \alpha^{(k+k)k} & \dots & \alpha^{(n-k)(n-1)} \end{pmatrix}$$

phd $\rightarrow 2 \quad (RS \text{ Cn,2})$

$$2 = t = \left\lceil \frac{n+1}{2} \right\rceil \rightarrow n-k = 4$$

$$d_{min} = n-k+1$$

Kodiertreter

mit eingesetztem : C

g. pun

g. pun $n = q-1$

g	n	k
2	1	-
3	2	-
5	4	0
7	6	2
11	10	6
13	12	8

$C_{RS}(6,2)$